

Script generated by TTT

Title: Petter: Programmiersprachen (20.01.2016)

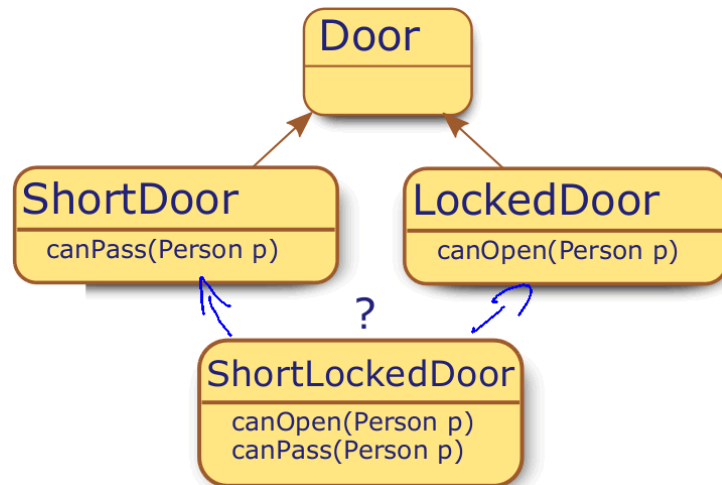
Date: Wed Jan 20 15:15:18 CET 2016

Duration: 45:47 min

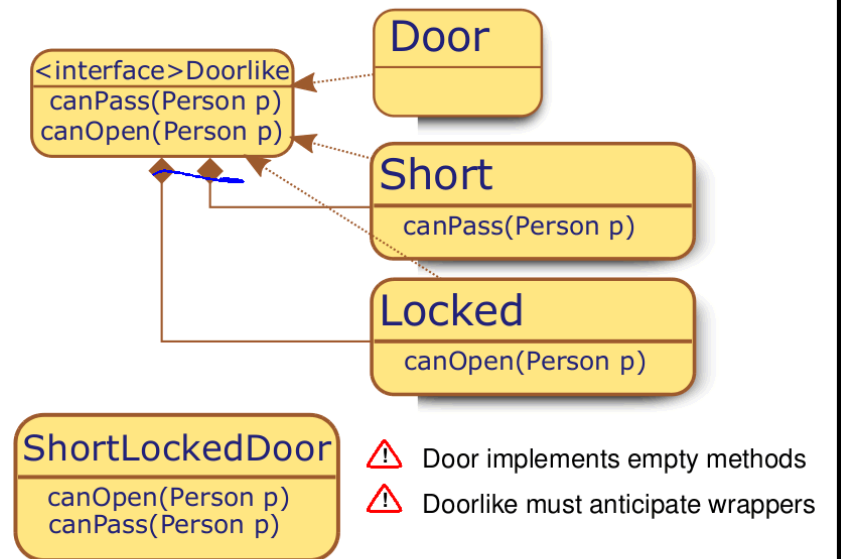
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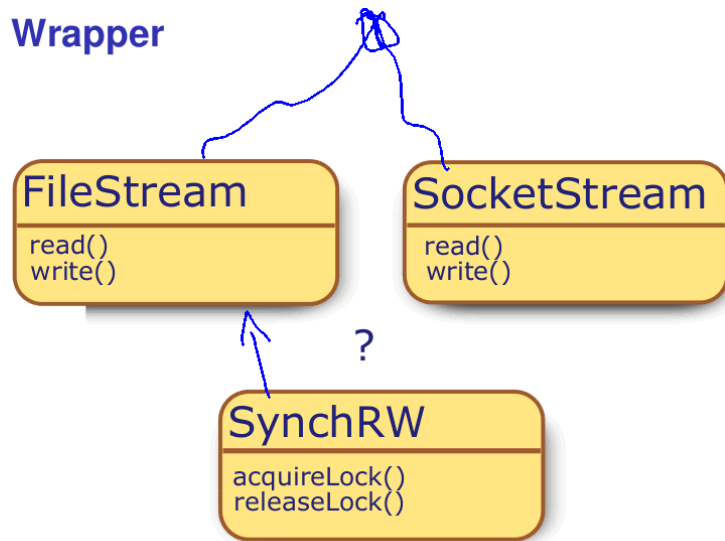
“What advanced techniques are there besides multiple implementation inheritance?”

The Adventure Game

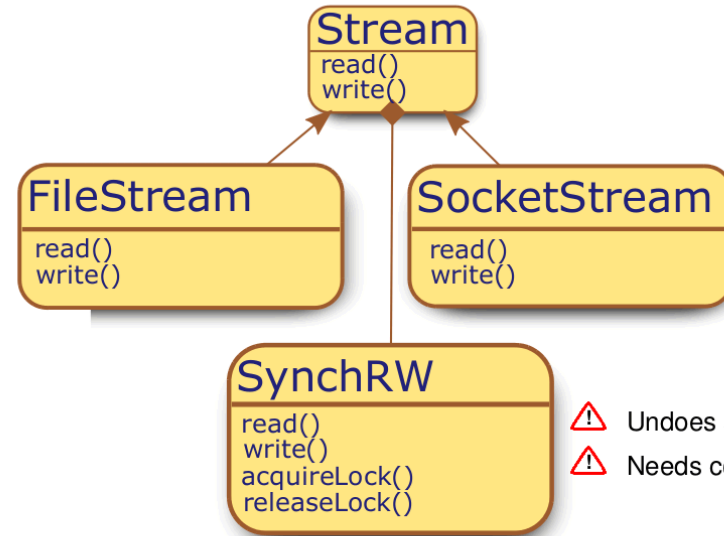


The Adventure Game





- ⚠ Cannot inherit from both separately
- ⚠ Creating new wrapping Classes duplicates code



- ⚠ Undoes specialization
- ⚠ Needs common ancestor

“Let’s go back to the basics of inheritance”

Classes and Methods

The building blocks for classes are

- a countable set of method *names* \mathcal{N}
- a countable set of method *bodies* \mathbb{B}

Classes map names to elements from the *flat lattice* \mathcal{B} (called bindings), consisting of:

- method bodies $\in \mathbb{B}$ or classes $\in \mathcal{C}$
- attribute offsets $\in \mathbb{N}^+$
- \perp (yet) undefined
- \top in conflict

and the partial order $\perp \sqsubseteq m \sqsubseteq \top$ for each $m \in \mathcal{B}$

Definition (Abstract Class $\in \mathcal{C}$)

A partial function $c : \mathcal{N} \mapsto \mathcal{B}$ is called abstract class.

Definition (Interface and Class)

An abstract class c is called (with pre being the preimage)

- interface* iff $\forall_{n \in \text{pre}(c)} \cdot c(n) = \perp$.
- (concrete) class* iff $\forall_{n \in \text{pre}(c)} \cdot \perp \sqsubseteq c(n) \sqsubseteq \top$.

Computing with Classes and Methods



Definition (Family of classes \mathcal{C})

We call the set of all maps from names to bindings the family of abstract classes $\mathcal{C} := \mathcal{N} \mapsto \mathcal{B}$.

Several possibilities for composing maps $\mathcal{C} \square \mathcal{C}$:

- the symmetric join \sqcup , defined componentwise:

$$(c_1 \sqcup c_2)(n) = b_1 \sqcup b_2 = \begin{cases} b_2 & \text{if } b_1 = \perp \text{ or } n \notin \text{pre}(c_1) \\ b_1 & \text{if } b_2 = \perp \text{ or } n \notin \text{pre}(c_2) \\ b_2 & \text{if } b_1 = b_2 \\ \top & \text{otherwise} \end{cases} \quad \text{where } b_i = c_i(n)$$

- in contrast, the asymmetric join $\sqcup\!\!\sqcup$ defined componentwise:

$$(c_1 \sqcup\!\!\sqcup c_2)(n) = \begin{cases} c_1(n) & \text{if } n \in \text{pre}(c_1) \\ c_2(n) & \text{otherwise} \end{cases}$$

Example: Smalltalk-Inheritance



Smalltalk inheritance

- is the archetype for inheritance in mainstream languages like Java or C#
- inheriting smalltalk-style establishes a reference to the parent

Definition (Smalltalk inheritance \triangleright)

Smalltalk inheritance is the binary operator $\triangleright : \mathcal{C} \times \mathcal{C} \mapsto \mathcal{C}$, defined by $c_1 \triangleright c_2 = \{\text{super} \mapsto c_2\} \sqcup\!\!\sqcup (c_1 \sqcup c_2)$

Example: Doors

$$\begin{aligned} \text{Door} &= \{\text{canPass} \mapsto \perp, \text{canOpen} \mapsto \perp\} \\ \text{LockedDoor} &= \{\text{canOpen} \mapsto 0x4204711\} \triangleright \text{Door} \\ &= \{\text{super} \mapsto \text{Door}\} \sqcup\!\!\sqcup (\{\text{canOpen} \mapsto 0x4204711\} \sqcup \text{Door}) \\ &= \{\text{super} \mapsto \text{Door}, \text{canOpen} \mapsto 0x4204711, \text{canPass} \mapsto \perp\} \end{aligned}$$

Excursion: Beta-Inheritance



In *Beta*-style inheritance

- the design goal is to provide security from replacement of a method by a different method.
- methods in parents dominate methods in subclass
- the keyword `inner` explicitly delegates control to the subclass

Definition (Beta inheritance \triangleleft)

Beta inheritance is the binary operator $\triangleleft : \mathcal{C} \times \mathcal{C} \mapsto \mathcal{C}$, defined by

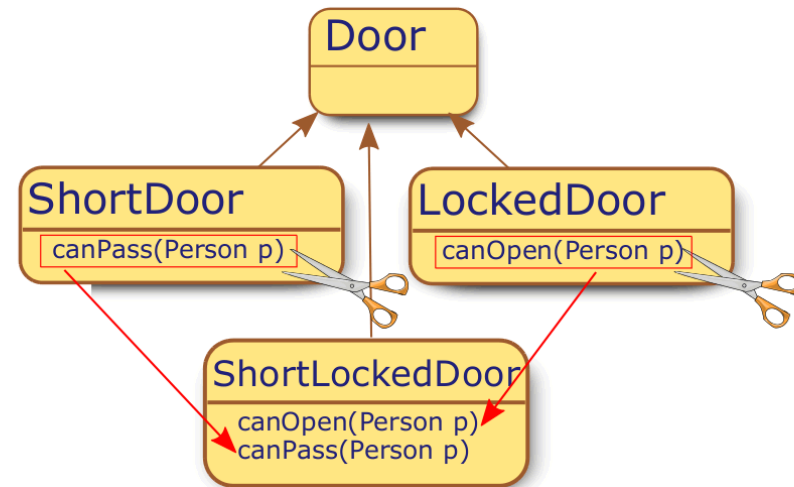
$$c_1 \triangleleft c_2 = \{\text{inner} \mapsto c_1\} \sqcup\!\!\sqcup (c_2 \sqcup c_1)$$

Example (equivalent syntax):

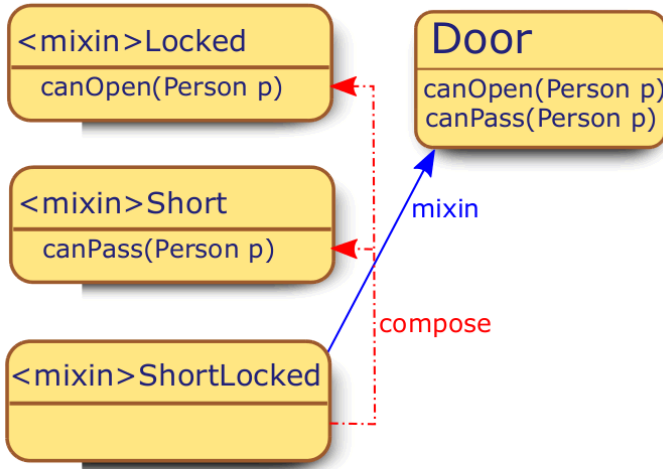
```
class Person {
  String name = "Axel Simon";
  public String toString(){ return name+inner.toString(); };
};

class Graduate extends Person {
  public extension String toString(){ return ", Ph.D."; };
};
```

Adventure Game with Code Duplication



Adventure Game with Mixins



Adventure Game with Mixins



```

class Door {
  boolean canOpen(Person p) { return true; };
  boolean canPass(Person p) { return p.size() < 210; };
}
mixin Locked {
  boolean canOpen(Person p){
    if (!p.hasItem(key)) return false; else return super.canOpen(p);
  }
}
mixin Short {
  boolean canPass(Person p){
    if (p.height()>1) return false; else return super.canPass(p);
  }
}
class ShortDoor = Short(Door);
class LockedDoor = Locked(Door);
mixin ShortLocked = Short o Locked;
class ShortLockedDoor = Short(Locked(Door));
class ShortLockedDoor2 = ShortLocked(Door);
    
```

Back to the blackboard!

Abstract model for Mixins



A Mixin is a *unary second order type expression*. In principle it is a curried version of the Smalltalk-style inheritance operator. In certain languages, programmers can create such mixin operators:

Definition (Mixin)

The mixin constructor $mixin : \mathcal{C} \mapsto (\mathcal{C} \mapsto \mathcal{C})$ is a unary class function, creating a unary class operator, defined by:

$$mixin(c) = \lambda x . c \triangleright x$$

⚠ Note: Mixins can also be composed \circ :

Example: Doors

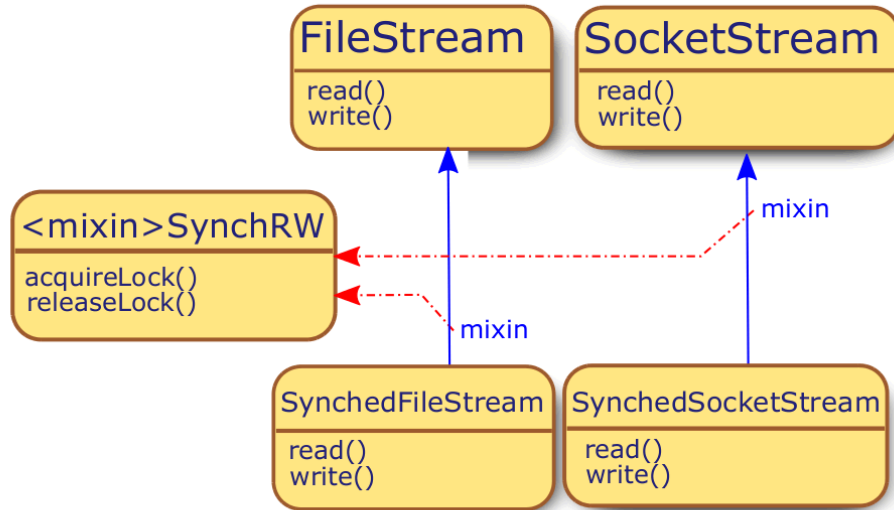
$$Locked = \{canOpen \mapsto 0x1234\}$$

$$Short = \{canPass \mapsto 0x4711\}$$

$$Composed = mixin(Short) \circ (mixin(Locked)) = \lambda x . Short \triangleright (Locked \triangleright x)$$

$$= \lambda x . \{super \mapsto Locked\} \uparrow (\{canOpen \mapsto 0x1234, canPass \mapsto 0x4711\} \triangleright x)$$

Wrapper with Mixins



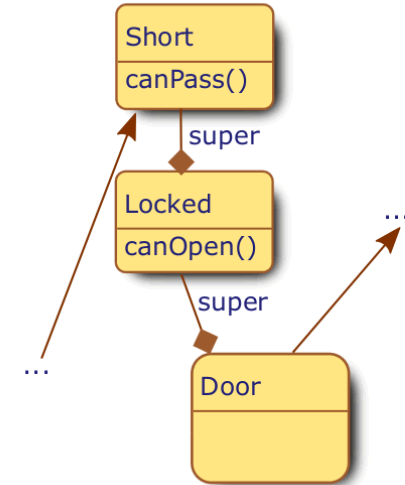
Mixins on Implementation Level



```

class Door {
    boolean canOpen(Person p)...
    boolean canPass(Person p)...
}
mixin Locked {
    boolean canOpen(Person p)...
}
mixin Short {
    boolean canPass(Person p)...
}
class ShortDoor
    = Short(Door);
class ShortLockedDoor
    = Short(Locked(Door));
...

ShortDoor d
    = new ShortLockedDoor();
    
```



⚠ non-static super-References
 ~> dynamic dispatching without precomputed virtual table

“Surely multiple inheritance is powerful enough to simulate mixins?”

Simulating Mixins in C++



```

template <class Super>
class SyncRW : public Super {
public: virtual int read(){
    acquireLock();
    int result = Super::read();
    releaseLock();
    return result;
};
virtual void write(int n){
    acquireLock();
    Super::write(n);
    releaseLock();
};
// ... acquireLock & releaseLock
};
    
```

Simulating Mixins in C++



```
template <class Super>
class LogOpenClose : public Super {
public: virtual void open(){
    Super::open();
    log("opened");
};
virtual void close(){
    Super::close();
    log("closed");
};
protected: virtual void log(char*s) { ... };
};
class MyDocument : public SyncRW, LogOpenClose<Document>> {};
```

True Mixins vs. C++ Mixins



True Mixins

- `super` natively supported
- Mixins as Template do not offer composite mixins
- C++ Type system not modular
- \rightsquigarrow Mixins have to stay source code
- Hassle-free simplified version of multiple inheritance

C++ Mixins

- Mixins reduced to templated superclasses
- Can be seen as coding pattern

Common properties of Mixins

- Linearization is necessary
- \rightsquigarrow Exact sequence of Mixins is relevant

Ruby



```
class Person
  attr_accessor :size
  def initialize
    @size = 160
  end
  def hasKey
    true
  end
end
class Door
  def canOpen (p)
    true
  end
  def canPass(person)
    person.size < 210
  end
end
```

```
module Short
  def canPass(p)
    p.size < 160 and super(p)
  end
end
module Locked
  def canOpen(p)
    p.hasKey() and super(p)
  end
end
class ShortLockedDoor < Door
  include Short
  include Locked
end
p = Person.new
d = ShortLockedDoor.new
puts d.canPass(p)
```

Ruby



```
class Door
  def canOpen (p)
    true
  end
  def canPass(person)
    person.size < 210
  end
end
module Short
  def canPass(p)
    p.size < 160 and super(p)
  end
end
module Locked
  def canOpen(p)
    p.hasKey() and super(p)
  end
end
```

```
module ShortLocked
  include Short
  include Locked
end
class Person
  attr_accessor :size
  def initialize
    @size = 160
  end
  def hasKey
    true
  end
end
p = Person.new
d = Door.new
d.extend ShortLocked
puts d.canPass(p)
```

Lessons Learned

- 1 Formalisms to model inheritance
- 2 Mixins provide soft multiple inheritance
- 3 Multiple inheritance can not compensate the lack of `super` reference
- 4 Full extent of mixins only when mixins are 1st class language citizens