Script generated by TTT

Title: Petter: Compiler Construction (11.06.2020)

- 31: ANSI C Example

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A Practial Example: Type Definitions in ANSI C

The C grammar distinguishes typename and identifier. Consider the following declarations:

```
typedef struct { int x,y } point_t
point_t origin;
```

A Practial Example: Type Definitions in ANSI C

A type definition is a *synonym* for a type expression. In C they are introduced using the *typedef* keyword. Type definitions are useful

as abbreviation:

```
typedef struct { int x; int y; } point_t;
• to construct recursive types:
```

Possible declaration in C:

more readable:

```
typedef struct list list_t;
struct list {
  int info;
  struct list* next;
}
struct list* head;

typedef struct list list_t;
struct list {
  int info;
  list_t* next;
}
list_t* head;
```

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A Practial Example: Type Definitions in ANSI C

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```
typedef struct { int x,y } point_t;
point_t origin;
```

Idea: in a parser action maintain a shared list between parser and scanner to communicate identifiers to report as typenames

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A Practial Example: Type Definitions in ANSI C

The C grammar distinguishes typename and identifier. Consider the following declarations:

```
typedef struct { int x,y } point_t;
point_t, origin;
```

Idea: in a parser action maintain a shared list between parser and scanner to communicate identifiers to report as typenames Relevant C grammar:

```
declaration
                       \rightarrow (declarationspecifier)^{+}declarator;
declarationspecifier → static | volatile · · · typedef
                             void | char | char · · · typename
declarator
                       → identifier | · · ·
```

Problem:

During reduction of the declaration, the scanner eagerly provides a new lookahead token, thus has already interpreted point_t in line 2 as identifier

```
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```

A Practial Example: Type Definitions in ANSI C: Solutions Relevant C grammar:

```
declaration
                          \rightarrow (declarationspecifier)^+declarator;
declaration specifier \rightarrow static \mid volatile \cdots typedef
                                 void | char | char · · · typename
declarator
                          \rightarrow identifier |\cdots|
```

Solution is difficult:

- \bullet try to fix the lookahead token class within the scanner-parser-channel \triangle a mess
- 2 add a rule to the grammar, to make it context-free:

```
typename
              \rightarrow
                     identifier
```

A Practial Example: Type Definitions in ANSI C: Solutions Relevant C grammar:

```
\rightarrow (declarationspecifier)^+declarator;
declaration
declaration specifier \rightarrow static \mid volatile \cdots typedef
                               | void | char | char · · · typename
declarator
                         → identifier | · · ·
```

Solution is difficult:

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A Practial Example: Type Definitions in ANSI C: Solutions Relevant C grammar:

```
declaration
                            \rightarrow (declarationspecifier)^+declarator:
declaration specifier \rightarrow \text{static} \mid \text{volatile} \cdots \text{typedef}
                                  | void | char | char · · · typename
declarator
                            → identifier | · · ·
```

Solution is difficult:

- try to fix the lookahead token class within the scanner-parser-channel \triangle a mess
- 2 add a rule to the grammar, to make it context-free:

```
typename
                                      identifier
Example input: (mytype1) (mytype2);
```

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A Practial Example: Type Definitions in ANSI C: Solutions Relevant C grammar:

```
\begin{array}{cccc} declaration & \rightarrow & (declarationspecifier)^+ declarator \,; \\ declarationspecifier & \rightarrow & \mathtt{static} \mid \mathtt{volatile} \cdots \mathtt{typedef} \\ & \mid \mathtt{void} \mid \mathtt{char} \mid \mathtt{char} \cdots \mathtt{typename} \\ declarator & \rightarrow & \mathtt{identifier} \mid \cdots \end{array}
```

Solution is difficult:

- 2 add a rule to the grammar, to make it context-free:

```
\begin{array}{ccc} & typename & \rightarrow & \mathtt{identifier} \\ & \texttt{Example input:} & (\texttt{mytype1}) \ (\texttt{mytype2}) \ ; \\ & castexpr & \rightarrow & (\mathtt{typename}) \ castexpr \\ & postfixexpr & \rightarrow & postfixexpr \ (expression) \end{array}
```

```
A Practial Example: Type Definitions in ANSI C: Solutions
    Relevant C grammar:
                                        \rightarrow (declarationspecifier)^+declarator;
                 declaration
                 declaration specifier \rightarrow \text{static} | \text{volatile} \cdots \text{typedef}
                                             | void | char | char · · · typename
                 declarator
                                         → identifier | · · ·
     Solution is difficult:
    ● try to fix the lookahead token class within the scanner-parser-channel 	 a mess
    2 add a rule to the grammar, to make it context-free:
                                  typename \rightarrow
                                                       identifier
        Example input:
                          (mytype1) (mytype2);
                                                                                 N=> E Tiad()/
                            castexpr \rightarrow (typename) castexpr
                           postfixexpr \rightarrow postfixexpr (expression)
    o register identifier as typename before lookahead is harmful
                 declaration \rightarrow (declarationspecifier)^+ declarator \{: act(); :\};
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```

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